

# ALGEBRAIC ORIGIN OF THE JONES POLYNOMIAL

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and  
Department of Mathematical Methods in Physics, Faculty of Physics,  
Warsaw University

October 22, 2010

# KNOTS AND LINKS

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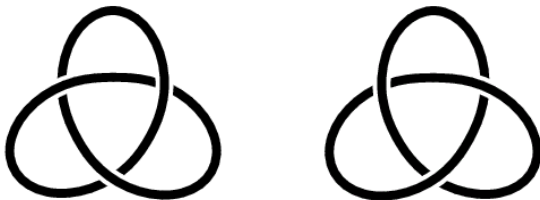
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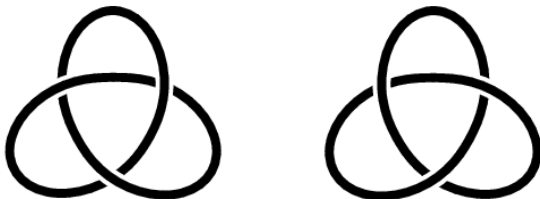
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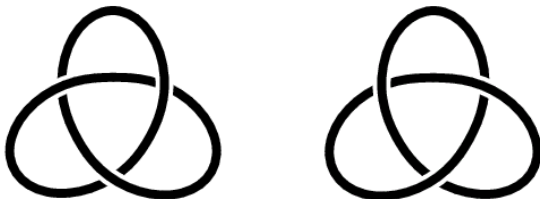
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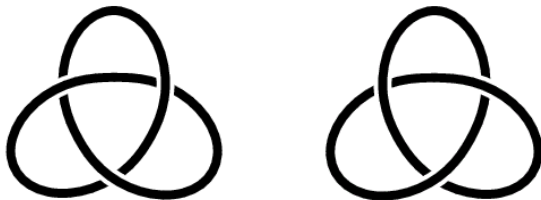
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- We are working in the PL-category.
- It is preferable to work with **oriented links** (each circle is oriented).



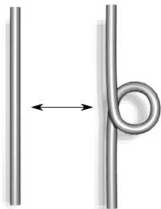
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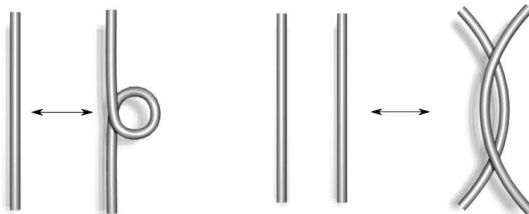
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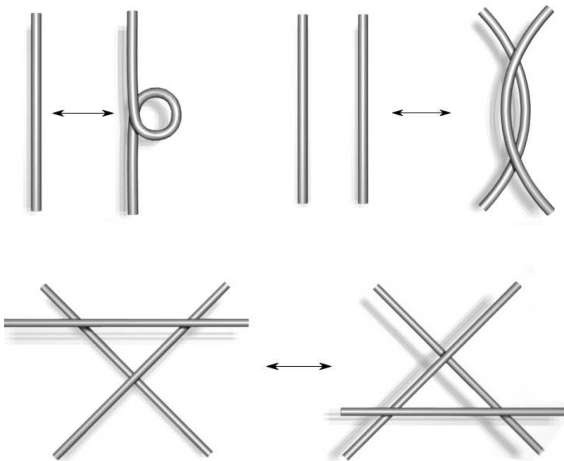
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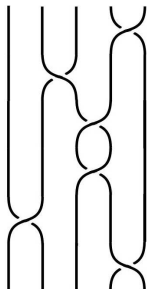
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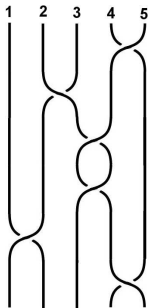
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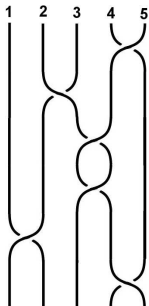
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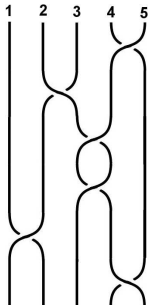
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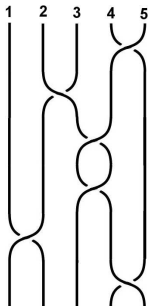
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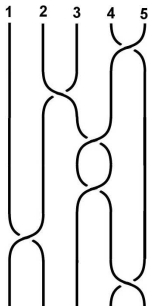
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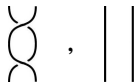
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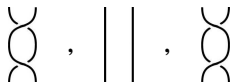
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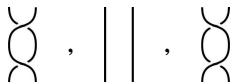
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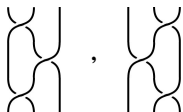


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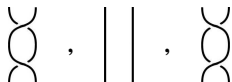


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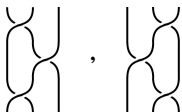


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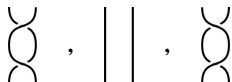
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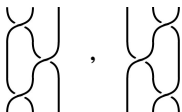
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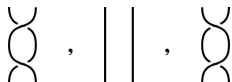
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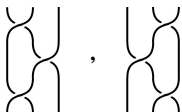
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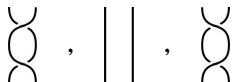
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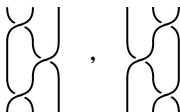
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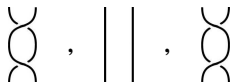
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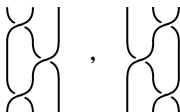
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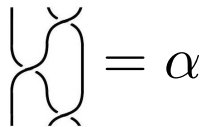
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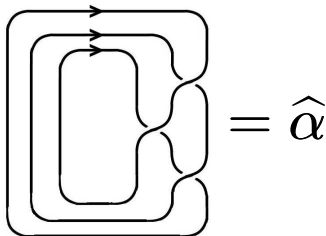
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$$\text{Diagram of a braid with two strands crossing twice} = \alpha$$

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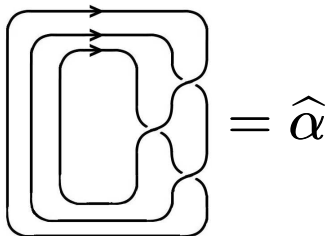
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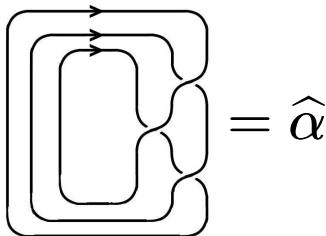
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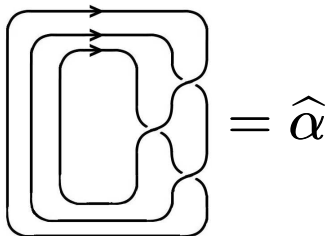
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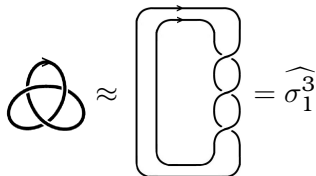
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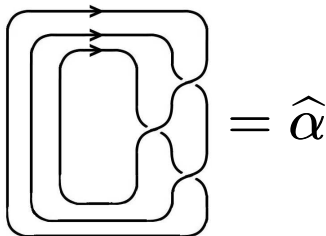
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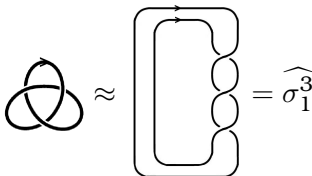
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(this is **Alexander's theorem**).

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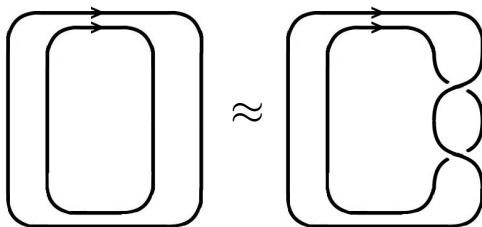
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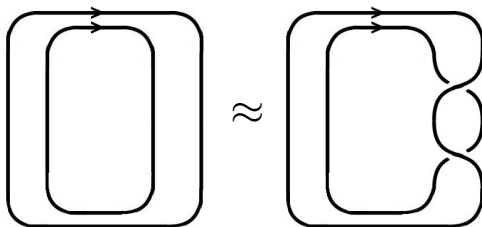
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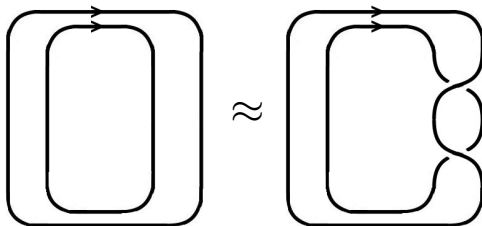
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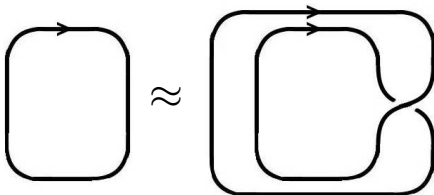
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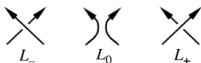
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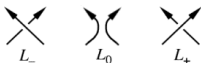
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- It follows that  $P_L(x, t)$  is uniquely determined by its value on the trivial knot:

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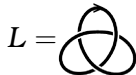
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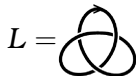
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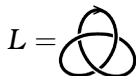
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